The proof of conjecture 36 from Ralf Stephan's "Prove or Disprove 100 Conjectures from the OEIS" found at http://front.math.ucdavis.edu/math.CO/0409509.

Conjecture 1 Let  $t(n) = |\phi(n) - n|$ . Then for any n > 0 we have that

$$t(n) = t(t(n) - n) \tag{1}$$

iff

$$n = 5 \cdot 2^k$$
 or  $n = 7 \cdot 2^k$  for some  $k > 0$ .

**Proof.** We first note that (1) is equivalent to

$$n - \phi(n) = t(-\phi(n))$$

$$n - \phi(n) = |\phi(-\phi(n)) + \phi(n)|$$

$$n - \phi(n) = \phi(\phi(n)) + \phi(n)$$

$$n = \phi(\phi(n)) + 2\phi(n)$$
(2)

Henceforth, we instead work to show that (2) is true only iff  $n = 5 \cdot 2^k$  or  $n = 7 \cdot 2^k$  for some k > 0.

If  $n=2^k5$ , with  $k\geq 1$  then  $\phi(n)=4\cdot 2^{k-1}=2^{k+1}$  and  $\phi(\phi(n))=2^k$  which clearly satisfies (2).

If  $n = 2^k 7$ , with  $k \ge 1$  then  $\phi(n) = 6 \cdot 2^{k-1} = 3 \cdot 2^k$  and  $\phi(\phi(n)) = 2 \cdot 2^{k-1} = 2^k$  which again satisfies (2).

We show that all other values for n fail.

- The cases  $1 \le n \le 6$  fail by inspection.
- If  $n \ge 7$ , then  $\phi(\phi(n)) + 2\phi(n)$  is even so (2) fails for all odd n.
- If  $n = 2^k$  for  $k \ge 2$ , then  $\phi(\phi(2^k)) + 2 \cdot \phi(2^k) = 2^{k-2} + 2 \cdot 2^{k-1} \ne 2^k$ . So  $n = 2^k$  fails for all  $k \ge 0$ .
- If  $n = 2^k M$  where M > 1 is odd and  $k \ge 1$ ,

$$n = \phi(\phi(n)) + 2\phi(n)$$

$$2^{k}M = \phi(\phi(2^{k}M)) + 2\phi(2^{k}M)$$

$$2^{k}M = \phi(2^{k-1} \cdot \phi(M)) + 2^{k} \cdot \phi(M)$$

$$2^{k}M = 2^{k-1} \cdot \phi(\phi(M)) + 2^{k} \cdot \phi(M) \text{ since } \phi(M) \text{ is even}$$

so we get the following further refinement of (2).

$$2M = \phi\left(\phi\left(M\right)\right) + 2\phi\left(M\right) \tag{3}$$

• If there exist two different odd primes p and q such that pq|M, then  $\phi(p)\phi(q)|\phi(M)$  so it follows that  $4|\phi(M)$ . In fact, we can write

$$\phi(p) \phi(q) = (p-1)(q-1) = 2^{\ell}N$$

for some  $\ell \geq 2$  and odd N. Therefore

$$\phi\left(2^{\ell}\right)\phi\left(N\right) = 2^{\ell-1}\phi\left(N\right)|\phi\left(\phi\left(M\right)\right).$$

If N=1, then  $p=2^{\alpha}+1$  and  $q=2^{\beta}+1$  for some  $\alpha \neq \beta$  so that  $\ell=\alpha+\beta\geq 3$ . If N>1 then  $\phi(N)$  is even. For both cases N=1 and N>1, we will have  $4|\phi(\phi(M))$ . So we conclude that (3) fails because 4 divides the RHS of (3) but not the LHS

Now we only need consider the case where M is an odd prime power..

• If  $n=2^k p$ , where  $k \geq 1$  and p is an odd prime, then (3) reduces to

$$2p = \phi(p-1) + 2(p-1)$$
  
 $2 = \phi(p-1)$ 

which fails for all cases other than p = 5 or 7.

• If  $n = 2^k p^j$ , where  $k \ge 1$ ,  $j \ge 2$  and p is an odd prime, then (3) reduces

$$\begin{array}{lll} 2^{k}p^{j} & = & \phi\left(\phi\left(2^{k}p^{j}\right)\right) + 2\phi\left(2^{k}p^{j}\right) \\ 2^{k}p^{j} & = & \phi\left(2^{k-1}\cdot(p-1)\,p^{j-1}\right) + 2\cdot2^{k-1}\,(p-1)\,p^{j-1} \\ 2^{k}p^{j} & = & 2^{k-1}\cdot\phi\left(p-1\right)\cdot(p-1)\,p^{j-2} + 2^{k}\,(p-1)\,p^{j-1} \text{ (since } p-1 \text{ is even and } (p-1,p) = 1) \\ 2p^{2} & = & \phi\left(p-1\right)\cdot(p-1) + 2\cdot\left(p^{2}-p\right) \\ 2p & = & \phi\left(p-1\right)\cdot(p-1) \end{array}$$

this fails for all odd p since p is clearly not a factor of the RHS.

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