

do!

BF's

MTAC 16 (1962)

1528

1530

~~2834~~

615

for calculation of Bessel
cal Bessel functions in digital
hev series," MTAC, v. IX, p.
xpansions of a class of hyper-
York, 1956, p. 57.

Majority Decision Functions of up to Six Variables

By S. Muroga, I. Toda, and M. Kondo

1. Introduction. Recently logical elements based essentially on the majority decision principle have been widely used in electronic computers. Among these elements are parametrons, magnetic cores, transistor-resistor logic, et cetera.

The logical behavior of such elements can be expressed by a model called a "majority decision element" with n Boolean inputs and one Boolean output, whose operation can be described in the form of a logical function called a "majority decision function".

This paper defines the canonical representative of each equivalence class in the classification of the majority decision functions by complementing and permuting variables and by complementing the output. Also, a method is proposed to obtain all the representatives with their optimum structures, and a table of the representatives of the majority decision functions of up to six variables is provided.

The reader should be familiar with the content of a previous paper by the authors, included as reference [1].

2. Majority Decision Functions. A "majority decision element" of n variables is a logical element with n Boolean inputs, x_1, x_2, \dots, x_n and one Boolean output. The output value of the element is

$$(1)* \quad \begin{aligned} &\text{one for } \sum_{i=1}^n w_i x_i \geq T \\ &\text{zero for } \sum_{i=1}^n w_i x_i \leq T - 1 \end{aligned}$$

where w_i is a prescribed constant real number called a "coupling weight" associated with the input x_i and T is also a prescribed constant real number called a "threshold."

In the case of parametrons or magnetic cores, the coupling weight w_i corresponds to the number of turns of the winding of the input x_i . The threshold T is related to the number of turns w_c for the constant input by the relation,

$$(2) \quad w_c = \sum_{i=1}^n w_i + 1 - 2T$$

where $w_c \geq 0$ means the constant of one is coupled to the element and $w_c < 0$ means the constant of zero.

A set of $(n + 1)$ real numbers $(w_1, w_2, \dots, w_n; T)$, which specifies the behavior of a majority decision element, will be called a "structure" of the element.

A logical function represented by a single majority decision element will be called a "majority decision function."

Received September 22, 1961.

* The term -1 on the right hand side is introduced as a normalizing factor of w_i 's and T .

For example, a majority decision element with the structure $(2, 1, 1; 2)$ represents a function $x_1 + x_2x_3$; hence, this function is a majority decision function. In contrast, the function $x_1x_2 + x_3x_4$ is not a majority decision function since it can not be realized by any single majority decision element.

3. Classification of Majority Decision Functions. Logical functions obtained from a given logical function f by the following operations are defined as equivalent functions with f :

- (1) Complementation of one or more input variables,
- (2) Permutation among input variables,
- (3) Complementation of f .

It is a well known fact that the logical functions can be classified into equivalent classes by this equivalent relation. Once a structure of a majority decision function is given, its equivalent functions can be easily realized in the same element by complementing and/or permuting input variables and/or by complementing the output. Thus, it is not necessary to determine the whole of the majority decision functions; it is sufficient to know the representatives of their equivalence classes. It should be noted that this limits the study to a much smaller number of functions.

In the case of general logical functions, it is difficult to extract systematically one representative from each equivalence class, but in the case of majority decision functions there is a way to define a canonical representative of each equivalence class from the intrinsic nature of majority decision functions.

The method of determining the canonical representative is described below. Hereafter in this section the majority decision function is assumed to have n non-vacuous variables.

Any majority decision function can be expressed by a polynomial without any complemented variable by appropriately complementing one or more variables (refer to [1], Section 3). Such a polynomial will be called a "positive polynomial." The way to complement the variables to obtain a positive polynomial from a given function is unique if complementing one variable more than once is prohibited. Hence we can restrict the possible representatives within positive polynomials. This is equivalent to the condition in which the representative should be realized by a majority decision element with positive coupling weights.

All the variables of a majority decision function can be ordered by a relation \gtrsim (refer to [1], Definition 3 and Theorem 1). Therefore, it is always possible for variables to be permuted and relabelled so that $x_1 \gtrsim x_2 \gtrsim \cdots \gtrsim x_m$ holds. This permutation can be uniquely determined except in the case of arbitrary permutations among some variables such as x_1, x_2, \dots, x_m for which $x_1 \sim x_2 \sim \cdots \sim x_m$ holds. But $x_1 \sim x_2 \sim \cdots \sim x_m$ means that the given function is symmetric with respect to these variables, and therefore the function is invariant under the permutations among x_1, x_2, \dots, x_m . Thus, the function for which $x_1 \gtrsim x_2 \gtrsim \cdots \gtrsim x_n$ holds is unique and can well be adopted as a possible representative. Of course, this is equivalent to the condition in which $w_1 \geq w_2 \geq \cdots \geq w_n$ holds for the representative majority decision element. Note that as a conclusion from the above requirements, we have $w_1 \geq w_2 \geq \cdots \geq w_n > 0$ except $w_n \leq 0$.

Only two functions left in each class satisfy both of the conditions just described.

If we denote one of the a majority decision function 2). A unique representative either of the two inequalit

Thus, it is shown that t
alent class of majority dec
ditions:

Conditions I.

- (1) A positive polynom
- (2) $x_1 \gtrsim x_2 \gtrsim \cdots \gtrsim$
- (3) f such that $f \subseteq f^*$.

Given a majority decis
sentative of the equivalen

**4. A Method to Obtain
Decision Functions.** From :
ming whether a given func
is possible, at least in princ
by applying the criterion t
ever, take an impractical
for large values of n , but th
reduced if we can confine t

Accordingly, a method
which includes all the repr
the criterion only to those
called "candidates" of the

Any positive majority c
 N , where M and N are bo
ables, x_2, x_3, \dots, x_n . Th
candidates within such fu
all the majority decision fu
here is one of the recursiv
spect to the number of va

Moreover, if we choose
I can be defined, then the
of the representatives of t

Then the restrictions in
Condition (1) will be t
struction.

Condition (2) requires

(3)

must hold for both M and
 N , it is necessary (Corolla

(4)

The structure $(2, 1, 1; 2)$ represents a majority decision function. In fact, it is a decision function since it can be represented by the following truth table:

Logical functions obtained
are defined as equivalent

.bles,

be classified into equivalent
a majority decision function
ed in the same element by
id/or by complementing the
hole of the majority decision
of their equivalence classes.
smaller number of functions.
ult to extract systematically
in the case of majority de-
representative of each equiv-
ision functions.

sent. π is described below.

by a polynomial without any
nting one or more variables
lled a "positive polynomial."
itive polynomial from a given
ore than once is prohibited.
within positive polynomials.
esentative should be realized
r weights.

n be ordered by a relation \succsim , it is always possible for $x_1 \succsim x_2 \succsim \cdots \succsim x_m$ holds. This is the case of arbitrary permutation which $x_1 \sim x_2 \sim \cdots \sim x_m$. A function is symmetric with respect to x_1, x_2, \dots, x_m if it is invariant under the permutations for which $x_1 \succsim x_2 \succsim \cdots \succsim x_n$ are representative. Of course, $w_1 \geq w_2 \geq \cdots \geq w_n$ holds for the w_i 's as a conclusion from the condition $w_i > 0$ except $w_n \leq 0$.

If we denote one of them by f , the other is the dual function f^* of f . But for a majority decision function, either $f^* \supseteq f$, or $f \supseteq f^*$ holds (refer to [1], Corollary 2). A unique representative of the equivalent class can be determined by requiring either of the two inequalities. If we adopt f such that $f \subseteq f^*$, this implies $w_i \leq 0$.

Thus, it is shown that there is a unique canonical representative in each equivalent class of majority decision functions which satisfies the following three conditions:

Conditions I

- (1) A positive polynomial,
 - (2) $x_1 \succsim x_2 \succsim \cdots \succsim x_n$,
 - (3) f such that $f \subseteq f^*$.

Given a majority decision function, we can now effectively obtain the representative of the equivalent class to which the given function belongs.

4. A Method to Obtain the Totality of the Representatives of the Majority Decision Functions. From Section 5 of [1] it can be determined by linear programming whether a given function is a majority decision function or not. Therefore, it is possible, at least in principle, to obtain the totality of majority decision functions by applying the criterion to all of 2^{2^n} logical functions of n variables. It will, however, take an impractically long time to solve 2^{2^n} linear programming problems for large values of n , but the length of time to perform computation will be greatly reduced if we can confine the scope of the functions to be tested.

Accordingly, a method is developed here to obtain a set of logical functions which includes all the representatives of majority decision functions and to apply the criterion only to those functions in the set. The functions in the set will be called "candidates" of the representatives.

Any positive majority decision function can be expressed in the form of $Mx_1 + N$, where M and N are both positive majority decision functions of $(n - 1)$ variables, x_2, x_3, \dots, x_n . Therefore, without loss of generality, we can restrict the candidates within such functions. This assumes that we have already obtained all the majority decision functions of $(n - 1)$ variables; hence the method described here is one of the recursive constructions of majority decision functions with respect to the number of variables.

Moreover, if we choose as the candidates those functions for which Conditions I can be defined, then the set of the candidates will certainly contain the totality of the representatives of the majority decision functions of n variables.

Then the restrictions imposed upon combinations of M and N will be examined.

Condition (1) will be trivially satisfied, for $Mx_1 + N$ is positive from its construction.

Condition (2) requires that the relation

$$(3) \quad x_3 \geq x_2 \geq \cdots \geq x_1$$

must hold for both M and N . Moreover, in order that $x_1 \gtrsim x_2$ may hold in $Mx_1 + N$, it is necessary (Corollary 1 of Reference [1]) that

(4) $m_1 \sqsupseteq m_2$

where

$$m_2 = M(0, x_3, \dots, x_n)$$

$$n_1 = N(1, x_3, \dots, x_n).$$

As the relation \gtrsim is an ordering relation (Theorem 1 of [1]), the relation

$$(5) \quad x_1 \gtrsim x_2 \gtrsim \dots \gtrsim x_n$$

follows from (3) and (4).

M and N are majority decision functions satisfying (3), hence the relations

$$(6) \quad m_1 \supseteq m_2 \text{ and } n_1 \supseteq n_2$$

where

$$m_1 = M(1, x_3, \dots, x_n)$$

$$n_2 = N(0, x_3, \dots, x_n)$$

hold (Corollary 1 of Reference [1]). From (4) and (6) we have

$$(7) \quad M \supset N.$$

From (3) in Conditions I, it is necessary that

$$(8) \quad f^* = N^*x_1 + M^*N^* \supseteq f = Mx_1 + N.$$

But as $M^*N^* = M^*$ from (7), (8) reduces to

$$(9) \quad M^* \supseteq N.$$

Thus, we choose as candidates those functions which satisfy the following conditions:

Conditions II

- (1) Both M and N are positive majority decision functions of $(n - 1)$ variables, x_2, x_3, \dots, x_n .
- (2) For both M and N , $x_2 \gtrsim x_3 \gtrsim \dots \gtrsim x_n$.
- (3) $m_2 \supseteq n_1$.
- (4) $M^* \supseteq N$.

By taking all the combinations of M and N which satisfy Conditions II, we can obtain the set of candidates of the representatives of majority decision functions of n variables.

M and N must satisfy (1) and (2) of Conditions II. Such functions are either canonical representatives of majority decision functions or their dual functions. Therefore, once the totality of representatives of majority decision functions of $(n - 1)$ variables are obtained, the scope within which functions M and N must be taken can be easily determined. In this way we can obtain the totality of the representatives of majority decision functions of n variables recursively.

The next problem is to examine each candidate to determine whether or not it is a majority decision function. If so, it is clearly a canonical representative of an equivalent class defined in the preceding section. The discrimination of majority decision functions from other functions can be accomplished by linear programming. The details will be found in Section 5 of [1].

5. Majority Decision Functions

described in Section 4, a particular MUSASINO-I, and all the variables were obtained.

The canonical representation Muroga [3] at that time, was completely.

The canonical representation in Table 1. The function $V = \sum_{i=1}^n w_i$, which is extent. Functions are express scripts. For instance, 12 +

In the same entry of the The optimum structure is namely, a structure which in [1]).

To establish the threshold element with a winding merely reversing the constant of one to the same

The numbers in this This is because f and f^* are in this paper and that in Table 1, the variables are shown, while the

By computing the number of majority decisions

6. Remarks on the Representative of Majority Functions

First, it is remarkable that is, Conditions II are satisfied by a single majority decision function.

Second, it is interesting that all integer-valued in such a solution of a system of equations.

A structure of a majority decision functions (Section 5 of [1]).

$$(10)$$

The third remark on the inequalities. It has been noted

, $x_n)$
, $x_n)$.
rem 1 of [1]), the relation
 x_n

satisfying (3), hence the relations
 $\supseteq n_2$
 \dots
, $x_n)$
 $x_n)$
and (6) we have
 $Mx_1 + N$.

ons which satisfy the following
cision functions of $(n - 1)$ var-
which satisfy Conditions II, we
tives of majority decision func-
ns II. Such functions are either
tions or their dual functions.
majority decision functions of
which functions M and N must
can obtain the totality of the
variables recursively.
e to determine whether or not
y a canonical representative of
The discrimination of majority
complished by linear program-
].

5. Majority Decision Functions of up to Six Variables. Following the procedure described in Section 4, a program was written for the parametron digital computer MUSASINO-I, and all the canonical representatives of the functions of up to six variables were obtained.

The canonical representatives of up to five variables had been obtained by S. Muroga [3] at that time, using a combinatorial method. Both results agreed completely.

The canonical representatives of the functions of up to six variables are shown in Table 1. The functions are numbered according to the magnitude of $V = \sum_{i=1}^n w_i$, which is expected to denote the complexity of functions to some extent. Functions are expressed by denoting the variables by means of their subscripts. For instance, $12 + 13 + 23$ stands for the function $x_1x_2 + x_1x_3 + x_2x_3$.

In the same entry of the table an optimum structure of the function is shown. The optimum structure is one with a minimum number of total turns of windings, namely, a structure which minimizes $(w_1 + w_2 + \dots + w_n + |w_e|)$ (Section 5 in [1]).

To establish the threshold T , the constant input of zero must be coupled to the element with a winding of $2T - V - 1$ turns. Dual functions can be realized by merely reversing the polarity of the constant input, that is, by coupling the constant of one to the same winding.

The numbers in this table are somewhat different from those shown in [1]. This is because f and f^* are considered to belong to the same equivalence class in this paper and that in Table 1 the numbers of functions of n (nonvacuous) variables are shown, while the numbers for up to n variables are shown in [1].

By computing the number of the members of each equivalent class, the total numbers of majority decision functions are obtained and shown in Table 2.

6. Remarks on the Results. Some remarks are added here concerning the representatives of majority decision functions of up to six variables.

First, it is remarkable that all the candidates proved to be true representatives, that is, Conditions II are sufficient for a function of up to six variables to be realized by a single majority decision element.

Second, it is interesting to note that the optimum structures (w_1, w_2, \dots, w_n) are all integer-valued in spite of the fact that the optimum structure is obtained as a solution of a system of inequalities of the form of equation (1).

A structure of a majority decision function is a solution of a system of 2^n linear inequalities (Section 5 of Reference [1]).

$$(10) \quad \begin{aligned} a_i > b_i & \quad : \quad \left\{ \begin{array}{l} a_i, i \mid 1, 2, \dots, 2^n \\ i \rightarrow 1, 2, \dots, n \end{array} \right\} \\ x = \begin{bmatrix} w_1 \\ w_2 \\ \vdots \\ w_n \\ T \end{bmatrix}, & \end{aligned}$$

The third remark concerns the structure of the solution space of these inequalities. It has been noted that for a majority decision function of up to five

TABLE 1
Representative Functions of Majority Decision Functions of up to Six Variables

No.	V	w_i	T	Representative Function	No.	V	w_i	T	Representative Function
$n = 2$									
$n = 3$									
1	2	11	2	12	19	9	32211	7	$123 + 1245 + 1345$
2	3	111	2	$12 + 13 + 23$	20	9	33111	6	$12 + 1345 + 2345$
3	4	211	3	123	21	9	33111	7	$123 + 124 + 125$
$n = 4$									
1	4	1111	3	$123 + 124 + 134 + 234$	22	9	42111	6	$12 + 134 + 135 + 145$
2	4	1111	4	1234	23	10	32221	6	$123 + 124 + 134 + 234 + 125 + 135 + 145$
3	5	2111	3	$12 + 13 + 14 + 234$	24	10	32221	7	$123 + 124 + 134 + 234$
4	5	2111	4	$123 + 124 + 134$	25	10	33211	6	$12 + 134 + 234 + 135 + 235$
5	6	2211	4	$12 + 134 + 234$	26	10	33211	8	$123 + 1245$
6	6	2211	5	$123 + 124$	27	10	42211	6	$12 + 13 + 145 + 2345$
7	6	3111	4	$12 + 13 + 14$	28	10	42211	7	$123 + 124 + 134 + 125 + 135$
8	7	3211	5	$12 + 134$	29	10	43111	7	$12 + 1345$
9	8	3221	5	$12 + 13 + 234$	30	11	33221	7	$123 + 124 + 134 + 234 + 125$
$n = 5$									
1	5	11111	3	$123 + 124 + 134 + 234 + 125 + 135 + 235 + 145 + 245 + 345$	31	11	33221	8	$123 + 124 + 1345 + 2345$
2	5	11111	4	$1234 + 1235 + 1245 + 1345 + 2345$	32	11	43211	7	$12 + 134 + 135 + 2345$
3	5	11111	5	12345	33	11	52211	7	$12 + 13 + 145$
4	6	21111	4	$123 + 124 + 134 + 125 + 135 + 145 + 2345$	34	12	33222	7	$123 + 124 + 134 + 234 + 125 + 135 + 235 + 145 + 245$
5	6	21111	5	$1234 + 1235 + 1245 + 1345$	35	12	43221	7	$12 + 134 + 234 + 135 + 145$
6	7	22111	4	$12 + 134 + 234 + 135 + 235 + 145 + 245$	36	12	43221	8	$123 + 124 + 134 + 125 + 2345$
7	7	22111	5	$123 + 124 + 125 + 1345 + 2345$	37	12	43221	9	$123 + 124 + 1345$
8	7	22111	6	$1234 + 1235 + 1245$	38	12	43311	7	$12 + 13 + 234 + 235$
9	7	31111	4	$12 + 13 + 14 + 15 + 2345$	39	12	52221	7	$12 + 13 + 14 + 2345$
10	7	31111	5	$123 + 124 + 134 + 125 + 135 + 145$	40	12	53211	8	$12 + 134 + 135$
11	8	22211	5	$123 + 124 + 134 + 234 + 125 + 135 + 235$	41	13	43321	8	$123 + 124 + 125 + 135 + 134 + 234$
12	8	22211	6	$123 + 1245 + 1345 + 2345$	42	13	53221	9	$123 + 124 + 125 + 134$
13	8	22211	7	$1234 + 1235$	43	13	53311	8	$12 + 13 + 2345$
14	8	32111	5	$12 + 134 + 135 + 145 + 2345$	44	14	43322	8	$123 + 124 + 125 + 134 + 135 + 145 + 234 + 235$
15	8	32111	6	$123 + 124 + 125 + 1345$	45	14	53321	8	$12 + 13 + 145 + 234$
16	8	41111	5	$12 + 13 + 14 + 15$	46	14	54221	9	$12 + 134 + 2345$
17	9	32211	5	$12 + 13 + 234 + 235 + 145$	47	15	54321	9	$12 + 134 + 135 + 234$
18	9	32211	6	$123 + 124 + 134 + 125 + 135 + 2345$	48	16	54322	9	$12 + 134 + 145 + 135 + 234 + 235$
$n = 6$									
No.	V	$w_1 \sim w_6$	T	Representative Function	No.	V	$w_1 \sim w_6$	T	Representative Function
$n = 6$									
5	7	211111	5	$1234 + 1235 + 1245 + 1345 + 1236 + 1246 + 1346 + 1256 + 1356 + 1456 + 23456$	47	11	521111	7	$12 + 134 + 135 + 145 + 136 + 146 + 156$
6	7	211111	6	$12345 + 12346 + 12356 + 12456 + 13456$	48	12	322221	7	$123 + 124 + 134 + 125 + 145 + 135 + 2345 + 2346 + 2356 + 2456 + 3456$
7	8	221111	5	$123 + 124 + 125 + 1345 + 126 + 1346 + 1356 + 1456 + 2345 + 2346 + 2356$	49	12	322221	8	$1234 + 1235 + 1245 + 1345 + 1236 + 1246 + 1346 + 1256 + 1356 + 1456 + 2345$
8	8	221111	6	$1234 + 1235 + 1245 + 1236 + 1246 + 1346 + 1356 + 1456 + 23456$	50	12	322221	9	$1234 + 1235 + 1245 + 1345 + 23456$
9	8	221111	7	$12345 + 12346 + 12356 + 12456$	51	12	322221	7	$123 + 124 + 134 + 125 + 126 + 1356 + 1456 + 234 + 2356 + 2456$
10	8	311111	5	$123 + 124 + 134 + 125 + 135 + 145 + 126 + 136 + 146 + 156 + 23456$	52	12	332211	8	$123 + 124 + 1345 + 1346 + 1256 + 2345 + 2346$
11	8	311111	6	$1234 + 1235 + 1245 + 1345 + 1236 + 1246 + 1346 + 1256 + 1356 + 1456$	53	12	332211	9	$1234 + 1235 + 1245 + 1236 + 1246 + 13456 + 23456$
12	9	222111	5	$123 + 124 + 131 + 125 + 135 + 126 + 136 + 1456 + 23456 + 2346 + 2356 + 2456$	54	12	332211	10	$1234 + 1235 + 1245 + 1236 + 1246 + 1256 + 13456$
				+ 3456	55	12	333111	7	$123 + 124 + 134 + 125 + 135 + 126 + 136 + 234 + 235 + 236$
					56	12	333111	9	$123 + 1245 + 13456 + 23456$

No.	V	$w_1 \sim w_6$	T	Representative Function	No.	V	$w_1 \sim w_6$	T	Representative Function
$n = 6$									
$n = 6$									
5	7	211111	5	$1234 + 1235 + 1245 + 1345 + 1236 + 1246 + 1346 + 1256 + 1356 + 1456 + 23456$	47	11	521111	7	$12 + 134 + 135 + 145 + 136 + 146 + 156$
6	7	211111	6	$12345 + 12346 + 12356 + 12456 + 13456$	48	12	322221	7	$123 + 124 + 134 + 125 + 145 + 135 + 2345 + 2346 + 2356 + 2456 + 3456$
7	8	221111	5	$123 + 124 + 125 + 1345 + 126 + 1346 + 1356 + 1456 + 2345 + 2346 + 2356$	49	12	322221	8	$1234 + 1235 + 1245 + 1345 + 1236 + 1246 + 1346 + 1256 + 1356 + 1456 + 2345$
8	8	221111	6	$1234 + 1235 + 1245 + 1236 + 1246 + 1346 + 1356 + 1456 + 23456$	50	12	322221	9	$1234 + 1235 + 1245 + 1345 + 23456$
9	8	221111	7	$12345 + 12346 + 12356 + 12456$	51	12	322221	7	$123 + 124 + 134 + 125 + 126 + 1356 + 1456 + 234 + 2356 + 2456$
10	8	311111	5	$123 + 124 + 134 + 125 + 135 + 145 + 126 + 136 + 146 + 156 + 23456$	52	12	332211	8	$123 + 124 + 1345 + 1346 + 1256 + 2345 + 2346$
11	8	311111	6	$1234 + 1235 + 1245 + 1345 + 1236 + 1246 + 1346 + 1256 + 1356 + 1456$	53	12	332211	9	$1234 + 1235 + 1245 + 1236 + 1246 + 13456 + 23456$
12	9	222111	5	$123 + 124 + 131 + 125 + 135 + 126 + 136 + 1456 + 23456 + 2346 + 2356 + 2456$	54	12	332211	10	$1234 + 1235 + 1245 + 1236 + 1246 + 1256 + 13456$
				+ 3456	55	12	333111	7	$123 + 124 + 134 + 125 + 135 + 126 + 136 + 234 + 235 + 236$
					56	12	333111	9	$123 + 1245 + 13456 + 23456$

No.	V	$w_1 \sim w_6$	T	Representative Function
2	5	11111	4	$1234 + 1235 + 1215 + 1345 + 2345$
3	5	11111	5	12345
4	6	21111	4	$123 + 124 + 131 + 125 + 135 + 145 + 2345$
5	6	21111	5	$1234 + 1235 + 1245 + 1345$
6	7	22111	4	$12 + 134 + 234 + 135 + 235 + 145 + 245$
7	7	22111	5	$123 + 124 + 125 + 1345 + 2345$
8	7	22111	6	$1234 + 1235 + 1245$
9	7	31111	4	$12 + 13 + 14 + 15 + 2345$
10	7	31111	5	$123 + 124 + 134 + 125 + 135 + 145$
11	8	22211	5	$123 + 124 + 134 + 234 + 125 + 135 + 235$
12	8	22211	6	$123 + 1245 + 1345 + 2345$
13	8	22211	7	$1234 + 1235$
14	8	32111	5	$12 + 134 + 135 + 145 + 2345$
15	8	32111	6	$123 + 124 + 125 + 1345$
16	8	41111	5	$12 + 13 + 14 + 15$
17	9	32211	5	$12 + 13 + 234 + 135 + 145$
18	9	32211	6	$123 + 124 + 134 + 125 + 135 + 2345$

No.	V	$w_1 \sim w_6$	T	Representative Function
1	6	111111	4	$1234 + 1235 + 1245 + 1345 + 1236 + 1246 + 1346 + 1256 + 1356 + 1456 + 2345$
2	6	111111	5	$12345 + 12346 + 12356 + 2456 + 3456$
3	6	111111	6	123456
4	7	211111	4	$123 + 124 + 134 + 125 + 135 + 145 + 2345 + 2346 + 2356 + 2456 + 3456$
5	7	211111	5	$1234 + 1235 + 1245 + 1345 + 1236 + 1246 + 1346 + 1256 + 1356 + 1456 + 2345$
6	7	211111	6	$12345 + 12346 + 12356 + 12456 + 13456$
7	8	221111	5	$123 + 124 + 125 + 1345 + 126 + 1346 + 1356 + 1456 + 2345 + 2346 + 2356 + 2456$
8	8	221111	6	$1234 + 1235 + 1245 + 1236 + 1246 + 1256 + 13456 + 23456$
9	8	221111	7	$12345 + 12346 + 12356 + 12456$
10	8	311111	5	$123 + 124 + 134 + 125 + 135 + 145 + 126 + 136 + 146 + 156 + 23456$
11	8	311111	6	$1234 + 1235 + 1245 + 1345 + 1236 + 1246 + 1346 + 1256 + 1356 + 1456$
12	9	222111	5	$123 + 124 + 134 + 125 + 135 + 126 + 136 + 1456 + 234 + 235 + 236 + 2456 + 3456$
13	9	222111	6	$123 + 1245 + 1345 + 1246 + 1346 + 1256 + 1356 + 2345 + 2346 + 2356$
14	9	222111	7	$1234 + 1235 + 1236 + 12456 + 13456 + 23456$
15	9	222111	8	$12345 + 12346 + 12356$
16	9	321111	5	$12 + 134 + 135 + 145 + 136 + 146 + 156 + 2345 + 2346 + 2356 + 2456$
17	9	321111	6	$123 + 124 + 125 + 1345 + 126 + 1346 + 1356 + 1456 + 23456$
18	9	321111	7	$1234 + 1235 + 1245 + 1236 + 1246 + 1256 + 13456$
19	9	411111	5	$12 + 13 + 14 + 15 + 16 + 23456$
20	9	411111	6	$123 + 124 + 134 + 125 + 135 + 145 + 126 + 136 + 146 + 156$
21	10	222211	6	$123 + 124 + 131 + 1256 + 1356 + 1456 + 234 + 2356 + 2456 + 3456$
22	10	222211	7	$1234 + 1235 + 1245 + 1345 + 1236 + 1246 + 1346 + 2345 + 2346 + 2356 + 2456$
23	10	222211	8	$1234 + 12356 + 12456 + 13456 + 23456$
24	10	222211	9	$12345 + 12346$
25	10	322111	6	$123 + 124 + 134 + 125 + 135 + 145 + 126 + 136 + 146 + 156 + 2345 + 2346 + 2356$
26	10	322111	7	$123 + 1245 + 1246 + 1346 + 1256 + 1356 + 23456$
27	10	322111	8	$1234 + 1235 + 1236 + 12456 + 13456$
28	10	331111	6	$12 + 1345 + 1346 + 1356 + 1456 + 2345 + 2346 + 2356$
29	10	331111	7	$123 + 124 + 125 + 126 + 13456 + 2345 + 2346 + 2356 + 2456$
30	10	331111	8	$1234 + 1235 + 1245 + 1236 + 1246 + 1256$
31	10	421111	6	$12 + 134 + 135 + 145 + 136 + 146 + 156 + 23456$
32	10	421111	7	$123 + 124 + 125 + 1345 + 126 + 1346 + 1356 + 1456$
33	10	511111	6	$12 + 13 + 14 + 15 + 16$
34	11	322211	6	$123 + 124 + 134 + 125 + 135 + 145 + 126 + 136 + 146 + 234 + 2356 + 2456 + 3456$
35	11	322211	7	$123 + 124 + 134 + 1256 + 1356 + 1456 + 2345 + 2346$
36	11	322211	8	$1234 + 1235 + 1245 + 1345 + 1236 + 1246 + 1346 + 23456$
37	11	322211	9	$1234 + 12356 + 12456 + 13456$
38	11	332111	6	$12 + 131 + 135 + 136 + 1456 + 234 + 235 + 236 + 2456$
39	11	332111	7	$123 + 124 + 125 + 1345 + 126 + 1346 + 1356 + 2345 + 2346 + 2356$
40	11	332111	8	$123 + 1245 + 1246 + 1256 + 13456 + 23456$
41	11	332111	9	$1234 + 1235 + 1236 + 12456$
42	11	422111	6	$12 + 13 + 145 + 146 + 156 + 2345 + 2346 + 2356$
43	11	422111	7	$123 + 124 + 134 + 125 + 135 + 126 + 136 + 1456 + 23456$
44	11	422111	8	$123 + 1245 + 1345 + 1246 + 1346 + 1256 + 1356$
45	11	431111	7	$12 + 1345 + 1346 + 1356 + 1456 + 23456$
46	11	431111	8	$123 + 124 + 125 + 126 + 13456$

2834

TABLE 2
The Number of Majority Decision Functions

615

<i>n</i>	Number of Logical Functions of up to <i>n</i> Variables	Number of Types of Logical Functions of <i>n</i> Variables*	Number of Types of Majority Decision Functions of <i>n</i> Variables	Number of Majority Decision Functions of <i>n</i> Variables	Number of Types of Self-Dual Majority Decision Functions of <i>n</i> Variables
1	4	1	1	2	1
2	16	2	1	8	0
3	256	10	3	72	1
4	65, 536	208	9	1, 536	1
5	4, 294, 967, 296	615, 904	48	86, 080	4
6	18, 446, 774, 073, 709, 551, 616	504	1528	14, 487, 040	14
			1530		

* These values are obtained from the results in References [4] and [5].

TABLE 3
The Maximum Values of Optimum Parameters of Majority Decision Functions

<i>n</i>	<i>w</i>	$V = \sum_{i=1}^n w_i$	<i>T</i>	<i>K</i>
2	1	2	2	3
3	2	4	3	5
4	3	8	5	9
5	5	16	9	17
6	9	33	18	35

variables the solution space of (10) is a pointed cone. That is, there is a certain point x_0 such that

$$(11) \quad Ax_0 \geq b$$

and any solution x of (10) can be written as

$$(12) \quad x = x_0 + x' \quad Ax' \geq 0.$$

This means the solution space of (10) is a cone with x_0 as a sole vertex. These structures for majority decision functions of six variables were examined and it was found that almost all the majority decision functions have solution space of a pointed cone but that 15 out of 504 representatives have spaces of non-cone structure. These functions are marked with * in Table 1.

Fourth, some maximum values of the optimum parameters are shown in Table 3, where V is the sum of coupling weights associated with input variables and K is the total number of turns of windings including the constant winding and the relation $K = 2T - 1$ holds. In Table 3, 26 functions have the maximum value 9 for a weight w and only one function attains the maximum value 33 of V ; there are 7 functions with maximum K of 35.

7. Acknowledgment. The authors wish to express their thanks to Mr. R. O. Winder, RCA Laboratories, Princeton, New Jersey, for his courtesy in comparing his data with ours, and to Dr. S. Takasu, Electrical Communication Laboratory, Tokyo, for his stimulating discussions.

International Business Machines Corporation
Thomas J. Watson Research Center
Yorktown Heights, New York

Electronics Research Section
Electrical Communication Laboratory
Musashino-shi, Tokyo, and

Electronics Research Section
Electrical Communication Laboratory
Musashino-shi, Tokyo

1. S. MUROGA, I. TODA, S. TAKASU, "Theory of majority decision elements," *J. Franklin Inst.*, v. 271, n. 5, May 1961, p. 376-418.
2. I. TODA, M. KONDO, S. MUROGA, "Majority decision functions of six variables," *Electrical Communication Laboratory Technical Journal*, v. 10, n. 3, 1961, p. 369-403, (in Japanese).
3. S. MUROGA, "A computer program to find Boolean functions representable by a single logical element based on a majority decision principle," *Electrical Communications Laboratory Technical Journal*, v. 8, n. 6, 1959, p. 614-622, (in Japanese).
4. D. SLEPIAN, "On the number of symmetry types of Boolean functions of n variables," *Canad. J. Math.*, v. 5, n. 2, 1953, p. 185-193.
5. B. ELSPAS, "Self-complementary types of Boolean functions," *IRE Trans. on Electronic Computers*, v. EC-9, n. 2, 1960, p. 264-266.
6. R. O. WINDER, "Single stage threshold logic," *AIEE Conference Paper* 60-1261, October 1960.
7. R. C. MINNICK, "Linear-input logic," *IRE Trans. on Electronic Computers*, v. EC-10, n. 1, 1961, p. 6-16.

TECHN

On the Number of Differences An

Summary. The computation of the number of differences requires numerical precision to obtain accuracy. One method of solution is to solve a difference equation over a limited range. With borderline stability, it is found that initial values eventually diverge.

To investigate this phenomenon arising in analytic number theory,

1. Introduction. The function $y(x)$ is equal to x and free of poles. It has been investigated by Buchstab, and de Bruijn has shown that

$$(1.1) \quad y(x) =$$

exists, where $y(x)$ is a function of x .

$$(1.2) \quad y(x) =$$

with $y(x) = 0$, $x < 0$.

The problem of computing $y(x)$ for $1 \leq x \leq 20$, was posed to the authors by Prof. H. W. Gould, and was described in [2], [3] which contains the differential equations, the formation concerning the function $y(x)$.

Tables of $y(x)$ are given for $1 \leq x \leq 5$ and of two or three digits, with additional effort, more significant digits can be obtained.

2. Computational Procedure

Received February 19, 1962