

New See/



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Dear Daniel,

Richard Stanley's problem in the February issue was too tantalizing to put down, so I enclose a solution.

Sincerely,

A handwritten signature in green ink, appearing to be "Don".

Donald E. Knuth  
Professor

DEK/pw

Enclosure: Solution to problem 10572

cc: Gessel, Sloane, Stanley, Wilf

**Solution to problem 10572.** The graphs in question can conveniently be called Stanley graphs. Let us say that a subspace *involves* coordinate  $k$  if it contains at least one vector such that the  $k$ th coordinate is nonzero; a subspace is called *full* if it involves all  $n$  coordinates. According to well-known principles of combinatorial enumeration [explained for example in *Combinatorial Species and Tree-like Structures* by Bergeron, Labelle, and Leroux (Cambridge University Press, 1997)],  $e^{-x} \sum g(n) x^n/n!$  is the exponential generating function for the number of full subspaces. Therefore it suffices to define a one-to-one correspondence between Stanley graphs on  $\{1, \dots, n\}$  and full subspaces of  $GF(2)^n$ .

Every subspace of dimension  $k$  has a canonical basis

$$u_1 = (u_{11}, \dots, u_{1n}), u_2 = (u_{21}, \dots, u_{2n}), \dots, u_k = (u_{k1}, \dots, u_{kn})$$

such that

$$u_{in_i} = 1, \quad u_{ij} = 0 \quad \text{for } j > n_i, \quad u_{ln_i} = 0 \quad \text{for } l \neq i,$$

for  $1 \leq i \leq k$ , where  $n \geq n_1 > n_2 > \dots > n_k \geq 1$ . For example, if  $n = 8$ ,  $k = 4$ ,  $n_1 = 8$ ,  $n_2 = 5$ ,  $n_3 = 3$ , and  $n_4 = 2$ , the basis has the schematic pattern

$$\begin{aligned} u_1 &= (u_{11}, 0, 0, u_{14}, u_{16}, u_{17}, 1), \\ u_2 &= (u_{21}, 0, 0, u_{24}, 1, 0, 0, 0), \\ u_3 &= (u_{31}, 0, 1, 0, 0, 0, 0, 0), \\ u_4 &= (u_{41}, 1, 0, 0, 0, 0, 0, 0). \end{aligned}$$

This basis, obtained by a familiar triangularization procedure, is uniquely defined by the subspace. [See D. E. Knuth, "Subspaces, subsets, and partitions," *J. Combinatorial Theory* **10** (1971), 178–180, for further elementary properties of the canonical basis.]

A subspace is full if and only if no component is zero in all of its canonical basis vectors. For example, the basis above defines a full subspace if and only if  $u_{11} \vee u_{21} \vee u_{31} \vee u_{41} = u_{14} \vee u_{24} = u_{16} = u_{17} = 1$ . We define a Stanley graph corresponding to each canonical basis by making vertex  $j$  adjacent to vertex  $n_i$  when  $u_{ij} = 1$ , for  $1 \leq i \leq k$ . For example, the subspace spanned by the vectors

$$\begin{aligned} u_1 &= (1, 0, 0, 0, 0, 1, 1, 1), \\ u_2 &= (0, 0, 0, 1, 1, 0, 0, 0), \\ u_3 &= (0, 0, 1, 0, 0, 0, 0, 0), \\ u_4 &= (1, 1, 0, 0, 0, 0, 0, 0) \end{aligned}$$

has edges  $\{1, 9\}$ ,  $\{7, 9\}$ ,  $\{8, 9\}$ ,  $\{4, 5\}$ , and  $\{1, 2\}$ . It is clear that this procedure defines a Stanley graph, since no vertex has neighbors both to the left and to the right. Furthermore,

the construction is reversible: A Stanley graph with exactly  $n - k$  vertices that have neighbors to the right corresponds to a full subspace of dimension  $k$ . Q.E.D.

**Appendix.** (The following remarks are mostly for my own records, but I include them here because I'm sending a copy of this letter to Ira Gessel, Richard Stanley, Herb Wilf, and Neil Sloane.) Stanley graphs are interesting combinatorial objects that do not seem to have been investigated before; at least, the number of such graphs

$$f(0), f(1), f(2), \dots = 1, 1, 2, 6, 158, 1330, 15414, 245578, 5382862, \dots$$

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is a sequence not yet in Sloane and Plouffe's *Encyclopedia of Integer Sequences*. Nor is the number of such graphs without isolated vertices,

$$1, 0, 1, 2, 11, 72, 677, 8686, 152191, 3632916, \dots$$

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obtained by multiplying the exponential generating function by  $e^{-x}$ ; this corresponds to full subspaces in which the canonical basis vectors all have at least two nonzero components. Nor is the number of *connected* Stanley graphs,

$$0, 1, 1, 2, 8, 52, 502, 6824, 127166, 3205924, \dots$$

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obtained from  $\ln \sum f(n)x^n/n!$ . Nor is the number of Stanley *trees*,

$$0, 1, 1, 2, 7, 36, 246, 2140, 21652, 260720, 3598120, \dots$$

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it is interesting to enumerate the latter (see below).

A graph on the vertices  $\{1, 2, \dots, n\}$  is a Stanley graph if and only if no vertex has both left and right neighbors. Thus every Stanley graph is bipartite: Every edge connects a vertex with right neighbors to a vertex with left neighbors. Conversely, the vertices of every bipartite graph can be labeled  $\{1, 2, \dots, n\}$  in such a way that we obtain a Stanley graph, simply by assigning the smallest labels to the vertices of one part and the largest labels to the vertices of the other.

Let  $\binom{n}{k}$  be the number of Stanley trees on  $n$  vertices such that  $k$  of the vertices have right neighbors and  $n - k$  have left neighbors. Let

$$\begin{aligned} s(x, y) &= \sum_{k,n} \binom{n}{k} \frac{x^k y^{n-k}}{n!} \\ &= \frac{xy}{2} + \frac{x^2y + xy^2}{6} + \frac{x^3y + 5x^2y^2 + xy^3}{24} + \dots \end{aligned}$$

Then by standard methods of combinatorial enumeration we find

$$y + \vartheta s(x, y) = y \exp(x + \vartheta_x s(x, y)),$$

where  $\vartheta_x$  is the differential operator  $x \frac{\partial}{\partial x}$  and  $\vartheta = x \frac{\partial}{\partial x} + y \frac{\partial}{\partial y}$ . It follows that we have the recurrence

$$\begin{bmatrix} n+1 \\ k \end{bmatrix} = \delta_{n1} \delta_{k1} + k \begin{bmatrix} n \\ k \end{bmatrix} + \begin{bmatrix} n \\ k-1 \end{bmatrix} + \sum_{m,j} \binom{n-1}{m-1} \begin{bmatrix} m \\ j \end{bmatrix} \begin{bmatrix} n+1-m \\ k-j \end{bmatrix} j;$$

these numbers  $\begin{bmatrix} n \\ k \end{bmatrix}$  form the "Gessel-Stanley triangle"

				0						
				0	0					
			0	1	0					
		0	1	1	0					
	0	1	5	1	0					
	0	1	17	17	1	0				
	0	1	49	146	49	1	0			
	0	1	129	922	922	129	1	0		
0	1	321	4887	11234	4887	321	1	0		

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By the correspondence above,  $\begin{bmatrix} n \\ k \end{bmatrix}$  is the number of ways to put 0s and 1s into a tableau shape (Ferrers diagram) with  $k$  rows and  $n - k$  columns in such a way that exactly  $n - 1$  entries are equal to 1, and there is no way to partition the rows and columns into two parts  $R_1 \cup R_2$  and  $C_1 \cup C_2$  such that the entries of  $R_1 \cap C_1$  and  $R_2 \cap C_2$  are entirely zero. For example, the 17 such arrays when  $n = 5$  and  $k = 2$  are

111	111	111	111	111	111	110	110	
1	10	100	01	010	001	101	011	
101	101	101	011	011	011	001	010	100
11	110	011	11	110	101	111	111	111.

We have  $\sum \begin{bmatrix} n \\ 2 \end{bmatrix} x^n = x^3/(1-x)(1-2x)^2$ , according to Sloane and Plouffe.

After writing the above, I happened to see Ira Gessel's paper in *Electronic J. Combinatorics* 3 (2) (1996), #R8 [*The Foata Festschrift*, 197-201], where the quantity  $\begin{bmatrix} n \\ k \end{bmatrix}$  is called  $u_{n-1, k-1, n-k-1}$ ; Gessel's  $u_{n,i,j}$  is the number of forests on vertices  $\{1, 2, \dots, n\}$  having  $i$  descents and  $j+1$  leaves. So Stanley trees have implicitly been considered before!